Deep Learning Sequence to Sequence models: Connectionist Temporal Classification

Sequence-to-sequence modelling

- Problem:
 - A sequence $X_1 \dots X_N$ goes in
 - A different sequence $Y_1 \dots Y_M$ comes out
- E.g.
 - Speech recognition: Speech goes in, a word sequence comes out
 - Alternately output may be phoneme or character sequence
 - Machine translation: Word sequence goes in, word sequence comes out
 - Dialog : User statement goes in, system response comes out
 - Question answering : Question comes in, answer goes out
- In general $N \neq M$
 - No synchrony between X and Y.



- Sequence goes in, sequence comes out
- No notion of "time synchrony" between input and output
 - May even not even maintain order of symbols
 - E.g. "I ate an apple" \rightarrow "Ich habe einen apfel gegessen"



- Or even seem related to the input
 - E.g. "My screen is blank" → "Can you check if your computer is plugged in?"

Case 1: Order-aligned but not time synchronous





- The input and output sequences happen in the same order
 - Although they may not be *time synchronous*, they can be "aligned" against one another
 - E.g. Speech recognition
 - The input speech can be aligned to the phoneme sequence output

Problems

- How do we perform *inference* on such a model Partially addressed
 - How to output time-asynchronous sequences
- How do we *train* such models

Problems

- How do we perform *inference* on such a model
 - How to output time-asynchronous sequences
- How do we *train* such models



Given the order-aligned output sequence with timing



- Given the order aligned output sequence with timing
 - Convert it to a time-synchronous alignment by repeating symbols
- Compute the divergence from the time-aligned sequence

$$DIV = \sum_{t} KL(Y_t, symbol_t) = -\sum_{t} \log Y(t, symbol_t)$$



$$DIV = \sum_{t} KL(Y_t, symbol_t) = -\sum_{t} \log Y(t, symbol_t)$$

• The gradient w.r.t the *t*-th output vector Y_t

$$\nabla_{Y_t} DIV = \begin{bmatrix} 0 & 0 & \dots & \frac{-1}{Y(t, symbol_t)} & 0 & \dots & 0 \end{bmatrix}$$

Zeros except at the component corresponding to the target aligned to that time



- Only the sequence of output symbols is provided for the training data
 - But no timing information



Iterative update: Problem

- Approach heavily dependent on initial alignment
- Prone to poor local optima
- Alternate solution: Do not commit to an alignment during any pass..

Recap: Training without alignment

- We know how to train if the alignment is provided
- Problem: Alignment is *not* provided
- Solution:
 - 1. Guess the alignment

2. Consider *all possible* alignments

Recap: The "aligned" table



The reason for suboptimality



- We *commit* to the single "best" estimated alignment
 - The most likely alignment

$$DIV = -\sum_{t} \log Y(t, symbol_{t}^{bestpath})$$

- This can be way off, particularly in early iterations, or if the model is poorly initialized
- Alternate view: there is a probability distribution over alignments of the target Symbol sequence (to the input)
 - Selecting a single alignment is the same as drawing a single sample from it
 - Selecting the most likely alignment is the same as deterministically always drawing the most probable value from the distribution

Averaging over all alignments



 Instead of only selecting the most likely alignment, use the statistical expectation over *all* possible alignments

$$DIV = E\left[-\sum_{t}\log Y(t, s_t)\right]$$

- Use the *entire distribution of alignments*
- This will mitigate the issue of suboptimal selection of alignment

The expectation over all alignments



$$DIV = E\left[-\sum_{t} \log Y(t, s_t)\right]$$

Using the linearity of expectation

$$DIV = -\sum_{t} E[\log Y(t, s_t)]$$

- This reduces to finding the expected divergence at each input

$$DIV = -\sum_{t} \sum_{S \in S_1 \dots S_K} P(s_t = S | \mathbf{S}, \mathbf{X}) \log Y(t, s_t = S)$$

17

The expectation over all alignments



The probability of aligning the specific symbol s at time t, given that unaligned sequence $S = S_0 \dots S_{K-1}$ and given the input sequence $X = X_0 \dots X_{N-1}$ We need to be able to compute this

$$DIV = -\sum_{t} E[\log Y(t, s_t)]$$

- This reduces to finding the expected divergence at each input

$$DIV = -\sum_{t} \sum_{S \in S_1 \dots S_K} P(s_t = S | \mathbf{S}, \mathbf{X}) \log Y(t, s_t = S)$$

18

A posteriori probabilities of symbols



$$P(s_t = S_r | \mathbf{S}, \mathbf{X}) \propto P(s_t = S_r, \mathbf{S} | \mathbf{X})$$

- $P(s_t = S_r, \mathbf{S} | \mathbf{X})$ is the total probability of all valid paths in the graph for target sequence \mathbf{S} that go through the symbol S_r (the r^{th} symbol in the sequence $S_0 \dots S_{K-1}$) at time t
- We will compute this using the "forward-backward" algorithm

A posteriori probabilities of symbols



•
$$P(s_t = S_r, \mathbf{S} | \mathbf{X})$$
 can be decomposed as

$$P(s_t = S_r, \mathbf{S} | \mathbf{X}) = P(S_0, \dots, S_r, \dots, S_{K-1}, s_t = S_r | \mathbf{X})$$

$$= P(S_0 \dots S_r, s_t = S_r, s_{t+1} \in succ(S_r), succ(S_r), \dots, S_{K-1} | \mathbf{X})$$
• Using Bayes Rule

$$= P(S_0 \dots S_r, s_t = S_r | \mathbf{X}) P(s_{t+1} \in succ(S_r), succ(S_r), \dots, S_{K-1} | S_0 \dots S_r, s_t = S_r \mathbf{X})$$

• The probability of the subgraph in the blue outline, times the conditional probability of the red-encircled subgraph, given the blue subgraph

Conditional independence y_0 y_1 $x = X_0 X_1 \dots X_{N-1}$ $H = H_0 H_1 \dots H_{N-1}$ y_1 \vdots

- **Dependency graph:** Input sequence $\mathbf{X} = X_0 X_1 \dots X_{N-1}$ governs hidden variables $\mathbf{H} = H_0 H_1 \dots H_{N-1}$
- Hidden variables govern output predictions y_0, y_1, \dots, y_{N-1} individually
- y_0, y_1, \dots, y_{N-1} are conditionally independent given **H**
- Since **H** is deterministically derived from **X**, y_0 , y_1 , ... y_{N-1} are also conditionally independent given **X**
 - This wouldn't be true if the relation between X and H were not deterministic or if X is unknown, or if the ys at any time went back into the net as inputs

A posteriori symbol probability



$$P(s_t = S_r, \mathbf{S} | \mathbf{X})$$

= $P(S_0 \dots S_r, s_t = S_r | \mathbf{X}) P(s_{t+1} \in succ(S_r), succ(S_r), \dots, S_{K-1} | \mathbf{X})$

- We will call the first term the *forward probability* $\alpha(t, r)$
- We will call the second term the *backward* probability $\beta(t, r)$

Computing $\alpha(t, r)$: Forward algorithm





 $\alpha(t,r) = P(S_0..S_r, s_t = S_r | \mathbf{X})$

 $\alpha(3, IY) = \alpha(2, B)y_3^{IY} + \alpha(2, IY)y_3^{IY}$

$$\alpha(t,r) = \sum_{q:S_q \in pred(S_r)} \alpha(t-1,q) Y_t^{S(r)}$$

• Where $pred(S_r)$ is any symbol that is permitted to come before an S_r and may include S_r

23

• q is its row index, and can take values r and r-1 in this example

Forward algorithm



 $\alpha(t,r) = \big(\alpha(t-1,r) + \alpha(t-1,r-1)\big)y_t^{S(r)}$

Forward algorithm



• Initialization:

$$\alpha(0,0) = y_0^{S(0)}, \quad \alpha(0,r) = 0, \ r > 0$$

• for $t = 1 \dots T - 1$

$$\alpha(t,0) = \alpha(t-1,0)y_t^{S(0)}$$

for $l = 1 \dots K - 1$
 $\alpha(t,l) = (\alpha(t-1,l) + \alpha(t-1,l-1))y_t^{S(l)}$

The final forward probability $\alpha(t,r)$



$$\alpha(T-1, K-1) = P(S_0 \dots S_{K-1} | \mathbf{X})$$

• The probability of the entire symbol sequence is the alpha at the bottom right node

SIMPLE FORWARD ALGORITHM

#N is the number of symbols in the target output
#S(i) is the ith symbol in target output
#y(t,i) is the output of the network for the ith symbol at time t
#T = length of input

```
#First create output table
For i = 1:N
    s(1:T,i) = y(1:T, S(i))
```

#The forward recursion

```
# First, at t = 1
alpha(1,1) = s(1,1)
alpha(1,2:N) = 0
for t = 2:T
    alpha(t,1) = alpha(t-1,1)*s(t,1)
    for i = 2:N
        alpha(t,i) = alpha(t-1,i-1) + alpha(t-1,i)
        alpha(t,i) *= s(t,i)
```

Can actually be done without explicitly composing the output table

Using 1...N and 1...T indexing, instead of 0...N-1, 0...T-1, for convenience of notation

SIMPLE FORWARD ALGORITHM

#N is the number of symbols in the target output
#S(i) is the ith symbol in target output
#y(t,i) is the network output for the ith symbol at time t
#T = length of input

```
#The forward recursion
```

```
# First, at t = 1
alpha(1,1) = y(1,S(1))
alpha(1,2:N) = 0
for t = 2:T
    alpha(t,1) = alpha(t-1,1)*y(t,S(1))
    for i = 2:N
        alpha(t,i) = alpha(t-1,i-1) + alpha(t-1,i)
        alpha(t,i) *= y(t,S(i))
```

Without explicitly composing the output table

Using 1..N and 1..T indexing, instead of 0..N-1, 0..T-1, for convenience of notation

A posteriori symbol probability



$$P(s_t = S_r, \mathbf{S} | \mathbf{X}) = \frac{\alpha(t, r)}{P(s_{t+1} \in succ(S_r), succ(S_r), \dots, S_{K-1} | \mathbf{X})}$$

- We will call the first term the *forward probability* $\alpha(t,r)$
- We will call the second term the *backward* probability $\beta(t,r)$



Bacward probability



 $\beta(t,r) = P(s_{t+1} \in succ(S_r), succ(S_r), \dots, S_{K-1} | \mathbf{X})$

• $\beta(t,r)$ is the probability of the exposed subgraph, not including the orange shaded box

Backward probability





Backward probability



 $y_{4}^{IY}\beta(4,1)$

 $\beta(3,1) = +$ $y_4^F \beta(4,2)$

Backward algorithm



Backward algorithm



$$\beta(t,r) = \sum_{q:S_q \in succ(S_r)} \beta(t+1,q) y_{t+1}^{S_q}$$

- The $\beta(t, r)$ is the total probability of the subgraph shown
- The $\beta(t, r)$ terms at any time t are defined recursively in terms of the $\beta(t + 1, q)$ terms at the next time

Backward algorithm



• Initialization:

$$\beta(T-1, K-1) = 1, \ \beta(T-1, r) = 0, \ r < K-1$$

• for t = T - 2 downto 0

$$\beta(t,K) = \beta(t+1,K)y_{t+1}^{S(K)}$$

for $r = K - 2 \dots 0$
• $\beta(t,r) = y_{t+1}^{S(l)}\beta(t+1,r) + y_{t+1}^{S(r+1)}\beta(t+1,r+1)$

SIMPLE BACKWARD ALGORITHM

#N is the number of symbols in the target output
#S(i) is the ith symbol in target output
#y(t,i) is the output of the network for the ith symbol at time t
#T = length of input

```
#First create output table
For i = 1:N
   s(1:T,i) = y(1:T, S(i))
```

#The backward recursion

```
# First, at t = T
beta(T,N) = 1
beta(T,1:N-1) = 0
for t = T-1 downto 1
    beta(t,N) = beta(t+1,N)*s(t+1,N)
    for i = N-1 downto 1
        beta(t,i) = beta(t+1,i)*s(t+1,i) + beta(t+1,i+1))*s(t+1,i+1)
```

Can actually be done without explicitly composing the output table

Using 1...N and 1...T indexing, instead of 0...N-1, 0...T-1, for convenience of notation
BACKWARD ALGORITHM

#N is the number of symbols in the target output
#S(i) is the ith symbol in target output
#y(t,i) is the output of the network for the ith symbol at time t
#T = length of input

#The backward recursion

Without explicitly composing the output table

Using 1...N and 1...T indexing, instead of 0...N-1, 0...T-1, for convenience of notation

The joint probability



$$P(s_t = S_r, \mathbf{S} | \mathbf{X}) = \frac{\alpha(t, r)\beta(t, r)}{\beta(t, r)}$$

- We will call the first term the *f* and provint $\alpha(t,r)$
- We will call the second term the packwar probability $\beta(t,r)$

Forward algo

Backward algo

The posterior probability



$$P(s_t = S_r, \mathbf{S} | \mathbf{X}) = \alpha(t, r)\beta(t, r)$$

• The *posterior* is given by

$$P(s_t = S_r | \mathbf{S}, \mathbf{X}) = \frac{P(s_t = S_r, \mathbf{S} | \mathbf{X})}{\sum_{s'_r} P(s_t = S'_r, \mathbf{S} | \mathbf{X})} = \frac{\alpha(t, r)\beta(t, r)}{\sum_{r', \alpha(t, r')} \beta(t, r')}$$

The posterior probability



• Let the posterior $P(s_t = S_r | \mathbf{S}, \mathbf{X})$ be represented by $\gamma(t, r)$

$$\gamma(t,r) = \frac{\alpha(t,r)\beta(t,r)}{\sum_{r'} \alpha(t,r')\beta(t,r')}$$

COMPUTING POSTERIORS

#N is the number of symbols in the target output
#S(i) is the ith symbol in target output
#y(t,i) is the output of the network for the ith symbol at time t
#T = length of input

#Assuming the forward are completed first alpha = forward(y, S) # forward probabilities computed beta = backward(y, S) # backward probabilities computed

#Now compute the posteriors

```
for t = 1:T
sumgamma(t) = 0
for i = 1:N
    gamma(t,i) = alpha(t,i) * beta(t,i)
    sumgamma(t) += gamma(t,i)
end
for i=1:N
    gamma(t,i) = gamma(t,i) / sumgamma(t)
```

Using 1..N and 1..T indexing, instead of 0..N-1, 0..T-1, for convenience of notation



$$DIV = -\sum_{t} \sum_{s \in S_0 \dots S_{K-1}} P(s_t = s | \mathbf{S}, \mathbf{X}) \log Y(t, s_t = s)$$

$$DIV = -\sum_{t} \sum_{r} \gamma(t, r) \log y_t^{S(r)}$$



$$DIV = -\sum_{t} \sum_{s \in S_0 \dots S_{K-1}} P(s_t = s | \mathbf{S}, \mathbf{X}) \log Y(t, s_t = s)$$
$$DIV = -\sum_{t} \sum_{r} \gamma(t, r) \log y_t^{\mathbf{S}(r)}$$

• The derivative of the divergence w.r.t the output Y_t of the net at any time:

$$\nabla_{Y_t} DIV = \begin{bmatrix} \frac{dDIV}{dy_t^{s_0}} & \frac{dDIV}{dy_t^{s_1}} \\ \frac{dV}{dy_t^{s_1}} & \frac{dDIV}{dy_t^{s_1}} \end{bmatrix} \dots \begin{bmatrix} \frac{dDIV}{dy_t^{s_1}} \\ \frac{dV}{dy_t^{s_1}} \\ \frac{dV}{dy_t^{s_1}} \end{bmatrix} \dots$$

- Components will be non-zero only for symbols that occur in the training instance



• The derivative of the divergence w.r.t the output Y_t of the net at any time:

$$\nabla_{Y_t} DIV = \begin{bmatrix} \frac{dDIV}{dy_t^{s_0}} & \frac{dDIV}{dy_t^{s_1}} \\ \frac{dV}{dy_t^{s_1}} & \cdots & \frac{dDIV}{dy_t^{s_{L-1}}} \end{bmatrix}$$

Components will be non-zero only for symbols that occur in the training instance



Components will be non-zero only for symbols that occur in the training instance



Derivative of the expected divergence



The derivatives at both these locations must be summed to get $\frac{dDIV}{dy_A^{IY}}$

$$DIV = -\sum_{t} \sum_{r} \gamma(t, r) \log y_{t}^{S(r)}$$

• The derivative of the divergence w.r.t any particular output of the network must sum over all instances of that symbol in the target sequence

$$\frac{dDIV}{dy_t^l} = -\frac{1}{y_t^l} \sum_{r:S(r)=l} \gamma(t,r)$$

- E.g. the derivative w.r.t y_t^{IY} will sum over both rows representing /IY/ in the above figure

COMPUTING DERIVATIVES

#N is the number of symbols in the target output
#S(i) is the ith symbol in target output
#y(t,i) is the output of the network for the ith symbol at time t
#T = length of input

#Assuming the forward are completed first
alpha = forward(y, S) # forward probabilities computed
beta = backward(y, S) # backward probabilities computed

```
# Compute posteriors from alpha and beta
gamma = computeposteriors(alpha, beta)
```

```
#Compute derivatives
```

```
for t = 1:T
  dy(t,1:L) = 0 # Initialize all derivatives at time t to 0
  for i = 1:N
     dy(t,S(i)) -= gamma(t,i) / y(t,S(i))
```

Using 1...N and 1...T indexing, instead of 0...N-1, 0...T-1, for convenience of notation

Overall training procedure for Seq2Seq case 1

/B/ /IY/ /F/ /IY/



 Problem: Given input and output sequences without alignment, train models

Overall training procedure for Seq2Seq case 1

• Step 1: Setup the network

- Typically many-layered LSTM



• Step 2: Initialize all parameters of the network

Overall Training: Forward pass

- Foreach training instance
 - **Step 3**: Forward pass. Pass the training instance through the network and obtain all symbol probabilities at each time



Overall training: Backward pass



- Foreach training instance
 - Step 3: Forward pass. Pass the training instance through the network and obtain all symbol probabilities at each time
 - Step 4: Construct the graph representing the specific symbol sequence in the instance. This may require having multiple rows of nodes with the same symbol scores

Overall training: Backward pass



- Foreach training instance:
 - Step 5: Perform the forward backward algorithm to compute $\alpha(t,r)$ and $\beta(t,r)$ at each time, for each row of nodes in the graph. Compute $\gamma(t,r)$.
 - **Step 6:** Compute derivative of divergence $\nabla_{Y_t} DIV$ for each Y_t

Overall training: Backward pass

- Foreach instance
 - Step 6: Compute derivative of divergence $\nabla_{Y_t} DIV$ for each Y_t

$$\nabla_{Y_t} DIV = \begin{bmatrix} \frac{dDIV}{dy_t^0} & \frac{dDIV}{dy_t^1} & \dots & \frac{dDIV}{dy_t^{L-1}} \end{bmatrix}$$
$$\frac{dDIV}{dy_t^l} = -\sum_{\substack{r:S(r)=l}} \frac{\gamma(t,r)}{y_t^l}$$

• Step 7: Backpropagate $\frac{dDIV}{dy_t^l}$ and aggregate derivatives over minibatch and update parameters

Story so far: CTC models

- Sequence-to-sequence networks which irregularly output symbols can be "decoded" by Viterbi decoding
 - Which assumes that a symbol is output at each time and *merges* adjacent symbols
- They require alignment of the output to the symbol sequence for training
 This alignment is generally not given
- Training can be performed by iteratively estimating the alignment by Viterbi-decoding and time-synchronous training
- Alternately, it can be performed by optimizing the expected error over *all* possible alignments
 - Posterior probabilities for the expectation can be computed using the forward backward algorithm

A key decoding problem

- Consider a problem where the output symbols are characters
- We have a decode: R R R E E E E D
- Is this the compressed symbol sequence RED or REED?

We've seen this before



• /G//F//F//IY//D/ or /G//F//IY//D/?

A key decoding problem

- We have a decode: R R R E E E E D
- Is this the symbol sequence RED or REED?
- Solution: Introduce an explicit extra symbol which serves to separate discrete versions of a symbol
 - A "blank" (represented by "-")
 - RRR---EE---DDD = RED
 - RR-E--EED = REED
 - RR-R---EE---D-DD = RREDD
 - R-R-R---E-EDD-DDDD-D = RRREEDDD
 - The next symbol at the end of a sequence of blanks is always a new character
 - When a symbol repeats, there must be at least one blank between the repetitions
- The symbol set recognized by the network must now include the extra blank symbol
 - Which too must be trained

• Note the extra "blank" at the output



• Note the extra "blank" at the output

<mark>/B/ /IY/ /F/ /IY/</mark>



Note the extra "blank" at the output

 X_0

 X_1

 X_2

 γ_4^b y_1^b γ_{5}^{b} y_2^b y_3^b y_6^b y_7^b y_0^{D} y_0^{AH} y_1^{AH} y_2^{AH} y_3^{AH} y_5^{AH} y_6^{AH} y_4^{AH} y_7^{AH} /AH/ y_0^B y_5^B y_1^B y_2^B y_3^B y_4^B y_6^B y_7^B /B/ y_3^D y_5^D y_0^D y_2^D y_6^D y_7^D y_1^D y_4^D /D/ y_0^{EH} y_1^{EH} y_6^{EH} y_2^{EH} y_3^{EH} y_4^{EH} y_5^{EH} y_7^{EH} y_8^{EH} /EH/ y_1^{IY} y_0^{IY} y_2^{IY} y_3^{IY} y_4^{IY} y_5^{IY} y_6^{IY} y_7^{IY} $y_8^{\prime\prime}$ /IY/ y_0^F y_1^F y_2^F y_3^F y_4^F y_5^F y_6^F y_7^F /F/ y_5^G y_0^G y_1^G y_2^G y_3^G y_4^G y_6^G y_7^G /G/

 X_4

 X_5

 X_6

 X_7

 X_3

/B/ /IY/ /F/ /IY/

 y_8^b

 y_8^{AH}

 y_8^B

 y_8^D

 y_8^F

 y_8^G

 X_8

Note the extra "blank" at the output

/B/ /IY/ /F/ /F/ /IY/





• The original method without blanks

 Changing the example to /B/ /IY/ /IY/ /F/ from /B/ /IY/ /F/ /IY/ for illustration

-	y_0^b	y_1^b	<i>y</i> ₂ ^{<i>b</i>}	y_3^b	<i>y</i> ^{<i>b</i>} ₄	y_5^b	y_6^b	y_7^b	<i>y</i> ₈ ^{<i>b</i>}
/B/	y_0^B	y_1^B	y_2^B	y_3^B	y_4^B	y_5^B	y_6^B	y_7^B	y_8^B
-	y_0^b	y_1^b	<i>y</i> ₂ ^{<i>b</i>}	y_3^b	<i>y</i> ^{<i>b</i>} ₄	y_5^b	y_6^b	<i>y</i> ₇ ^{<i>b</i>}	<i>y</i> ₈ ^{<i>b</i>}
/IY/	y_0^{IY}	y_1^{IY}	y_2^{IY}	y_3^{IY}	y_4^{IY}	y_5^{IY}	y_6^{IY}	y_7^{IY}	y_8^{IY}
_	y_0^b	y_1^b	<i>y</i> ₂ ^{<i>b</i>}	y_3^b	y_4^b	y_5^b	y_6^b	y_7^b	<i>y</i> ₈ ^{<i>b</i>}
/IY/	y_0^{IY}	y_1^{IY}	y_2^{IY}	y_3^{IY}	y_4^{IY}	y_5^{IY}	y_6^{IY}	y_7^{IY}	y_8^{IY}
-	y_0^b	y_1^b	<i>y</i> ₂ ^{<i>b</i>}	y_3^b	y_4^b	y_5^b	y_6^b	y_7^b	<i>y</i> ₈ ^{<i>b</i>}
/F/	y_0^F	y_1^F	y_2^F	y_3^F	y_4^F	y_5^F	y_6^F	y_7^F	y_8^F
-	y_0^b	y_1^b	y_2^b	y_3^b	y_4^b	y_5^b	y_6^b	y_7^b	y_8^b

- With blanks
- Note: a row of blanks between any two symbols
- Also blanks at the very beginning and the very end



 Add edges such that all paths from initial node(s) to final node(s) unambiguously represent the target symbol sequence



 The first and last column are allowed to also end at initial and final blanks



- The first and last column are allowed to also end at initial and final blanks
- Skips are permitted across a blank, but only if the symbols on either side are different
 - Because a blank is *mandatory between repetitions of a symbol* but *not required between distinct symbols*

Composing the graph

```
\#N is the number of symbols in the target output \#S(i) is the ith symbol in target output
```

```
#Compose an extended symbol sequence Sext from S, that has the blanks
#in the appropriate place
#Also keep track of whether an extended symbol Sext(j) is allowed to connect
#directly to Sext(j-2) (instead of only to Sext(j-1)) or not
```

```
function [Sext, skipconnect] = extendedsequencewithblanks(S)
    i = 1
    for i = 1:N
        Sext(j) = b' # blank
        skipconnect(j) = 0
        j = j+1
        Sext(j) = S(i)
        if (i > 1 \&\& S(i) != S(i-1))
            skipconnect(j) = 1
        else
            skipconnect(j) = 0
        i = i+1
    end
    Sext(j) = 'b'
    skipconnect(j) = 0
    return Sext, skipconnect
```

Using 1..N and 1..T indexing, instead of 0..N-1, 0..T-1, for convenience of notation

Example of using blanks for alignment: Viterbi alignment with blanks

MODIFIED VITERBI ALIGNMENT WITH BLANKS

```
[Sext, skipconnect] = extendedsequencewithblanks(S)
N = length(Sext)  # length of extended sequence
```

Backtrace

```
AlignedSymbol(T) = Bscr(T,N) > Bscr(T,N-1) ? N, N-1;
for t = T downto 1
    AlignedSymbol(t-1) = BP(t,AlignedSymbol(t))
```

Using 1...N and 1...T indexing, instead of 0...N-1, 0...T-1, for convenience of notation

Modified Forward Algorithm



• Initialization:

$$-\alpha(0,0) = y_0^b, \alpha(0,1) = y_0^b, \alpha(0,r) = 0 \quad r > 1$$

Modified Forward Algorithm y_4^b γ_2^b y_0^b v_1^b v_2^b y_5^b y_6^b y_7^b y_8^b y_3^B y_5^B y_6^B y_7^B y_8^B y_0^B y_1^B y_2^B y_4^B /B/ y_5^b y_7^b y_8^b y_2^b v_2^b y_4^b y_1^b y_6^b /IY/ y_1^{IY} y_8^{IY} y_2^{IY} y_4^{IY} y_5^{IY} y_6^{IY} y_7^{IY} y_3^{IY} y_8^b y_1^b y_2^b y_3^b y_4^b y_5^b y_6^b y_7^b _ y_1^{IY} y_2^{IY} y_3^{IY} y_4^{IY} y_5^{IY} y_6^{IY} y_7^{IY} y_8^{IY} /IY/ _ /F/ y_1^b y_2^b y_8^b y_3^b y_4^b y_5^b y_6^b y_7^b y_1^F y_2^F y_3^F y_4^F y_6^F y_8^F y_5^F y_7^F y_1^b y_3^b y_4^b y_8^b y_2^b y_6^b y_7^b t

• Iteration t = 1:N:

$$\alpha(t,r) = (\alpha(t-1,r) + \alpha(t-1,r-1))y_t^{S(r)}$$

• If $S(r) = "-"$ or $S(r) = S(r-2)$

 $\alpha(t,r) = (\alpha(t-1,r) + \alpha(t-1,r-1) + \alpha(t-1,r-2))y_t^{S(r)}$

• Otherwise

 $\alpha(t,r) = \sum \alpha(t-1,q) Y_t^{S(r)}$

 $q:S_a \in \overline{pred}(S_r)$

FORWARD ALGORITHM (with blanks)

[Sext, skipconnect] = extendedsequencewithblanks(S)
N = length(Sext) # Length of extended sequence

```
#The forward recursion
# First, at t = 1
alpha(1,1) = y(1,Sext(1)) #This is the blank
alpha(1,2) = y(1,Sext(2))
alpha(1,3:N) = 0
for t = 2:T
    alpha(t,1) = alpha(t-1,1)*y(t,Sext(1))
    for i = 2:N
        alpha(t,i) = alpha(t-1,i-1) + alpha(t-1,i))
        if (skipconnect(i))
            alpha(t,i) += alpha(t-1,i-2)
        alpha(t,i) *= y(t,Sext(i))
```

Without explicitly composing the output table

Using 1...N and 1...T indexing, instead of 0...N-1, 0...T-1, for convenience of notation
Modified Backward Algorithm



• Initialization:

$$\beta(T - 1, 2K - 1) = \beta(T - 1, 2K - 2) = 1$$

$$\beta(T - 1, r) = 0 \quad r < 2K - 2$$

Modified Backward Algorithm



• Iteration:

$$\beta(t,r) = \beta(t+1,r)y_{t+1}^{S(r)} + \beta(t+1,r+1)y_{t+1}^{S(r+1)}$$

• If S(r) = "-" or S(r) = S(r+2)

 $\beta(t,r) = \beta(t+1,r)y_{t+1}^{S(r)} + \beta(t+1,r+1)y_{t+1}^{S(r+1)} + \beta(t+1,r+2)y_{t+1}^{S(r+2)}$

• Otherwise

 $\beta(t,r) = \sum \qquad \beta(t+1,q)y_{t+1}^{S_q}$

 $q:S_a \in \overline{succ}(S_r)$

BACKWARD ALGORITHM WITH BLANKS

```
[Sext, skipconnect] = extendedsequencewithblanks(S)
N = length(Sext) # Length of extended sequence
```

#The backward recursion

```
# First, at t = T
beta(T,N) = 1
beta(T,N-1) = 1
beta(T,1:N-2) = 0
for t = T-1 downto 1
    beta(t,N) = beta(t+1,N)*y(t+1,Sext(N))
    for i = N-1 downto 1
        beta(t,i) = beta(t+1,i)*y(t+1,Sext(i)) + beta(t+1,i+1))*y(t+1,Sext(i+1))
        if (i<N-2 && skipconnect(i+2))
            beta(t,i) += beta(t+1,i+2)*y(t+1,Sext(i+2))
```

Without explicitly composing the output table

Using 1..N and 1..T indexing, instead of 0..N-1, 0..T-1, for convenience of notation

The rest of the computation

- Posteriors and derivatives are computed exactly as before
- But using the extended graphs with blanks

COMPUTING POSTERIORS

```
[Sext, skipconnect] = extendedsequencewithblanks(S)
N = length(Sext) # Length of extended sequence
```

```
#Assuming the forward are completed first
alpha = forward(y, Sext)  # forward probabilities computed
beta = backward(y, Sext)  # backward probabilities computed
```

```
#Now compute the posteriors
for t = 1:T
   sumgamma(t) = 0
   for i = 1:N
      gamma(t,i) = alpha(t,i) * beta(t,i)
      sumgamma(t) += gamma(t,i)
   end
   for i=1:N
      gamma(t,i) = gamma(t,i) / sumgamma(t)
```

Using 1...N and 1...T indexing, instead of 0...N-1, 0...T-1, for convenience of notation

COMPUTING DERIVATIVES

```
[Sext, skipconnect] = extendedsequencewithblanks(S)
N = length(Sext) # Length of extended sequence
```

#Assuming the forward are completed first
alpha = forward(y, Sext) # forward probabilities computed
beta = backward(y, Sext) # backward probabilities computed

```
# Compute posteriors from alpha and beta
gamma = computeposteriors(alpha, beta)
```

```
#Compute derivatives
for t = 1:T
    dy(t,1:L) = 0 #Initialize all derivatives at time t to 0
    for i = 1:N
        dy(t,Sext(i)) -= gamma(t,i) / y(t,Sext(i))
```

Using 1...N and 1...T indexing, instead of 0...N-1, 0...T-1, for convenience of notation

Overall training procedure for Seq2Seq with blanks

/B/ /IY/ /F/ /IY/



 Problem: Given input and output sequences without alignment, train models

Overall training procedure

Step 1: Setup the network
 Typically many-layered LSTM



Step 2: Initialize all parameters of the network
 – Include a "blank" symbol in vocabulary

Overall Training: Forward pass

- Foreach training instance
 - **Step 3**: Forward pass. Pass the training instance through the network and obtain all symbol probabilities at each time, including blanks



Overall training: Backward pass



- Foreach training instance
 - Step 3: Forward pass. Pass the training instance through the network and obtain all symbol probabilities at each time
 - Step 4: Construct the graph representing the specific symbol sequence in the instance. Use appropriate connections if blanks are included

Overall training: Backward pass



- Foreach training instance:
 - Step 5: Perform the forward backward algorithm to compute $\alpha(t,r)$ and $\beta(t,r)$ at each time, for each row of nodes in the graph using the modified forward-backward equations. Compute a posteriori probabilities $\gamma(t,r)$ from them
 - **Step 6:** Compute derivative of divergence $\nabla_{Y_t} DIV$ for each Y_t

Overall training: Backward pass

- Foreach instance
 - **Step 6:** Compute derivative of divergence $\nabla_{Y_t} DIV$ for each Y_t

$$\nabla_{Y_t} DIV = \begin{bmatrix} \frac{dDIV}{dy_t^0} & \frac{dDIV}{dy_t^1} & \dots & \frac{dDIV}{dy_t^{L-1}} \end{bmatrix}$$
$$\frac{dDIV}{dy_t^l} = -\sum_{\substack{r:S(r)=l}} \frac{\gamma(t,r)}{y_t^{S(r)}}$$

• Step 7: Backpropagate $\frac{dDIV}{dy_t^l}$ and aggregate derivatives over minibatch and update parameters

CTC: Connectionist Temporal Classification

- The overall framework we saw is referred to as CTC
- Applies to models that output order-aligned, but time-asynchronous outputs

Returning to an old problem: Decoding



- The greedy decode computes its output by finding the most likely symbol at each time and merging repetitions in the sequence
- This is in fact a *suboptimal* decode that actually finds the most likely *time-synchronous* output sequence
 - Which is not necessarily the most likely *order-synchronous* sequence

Greedy decodes are suboptimal

- Consider the following candidate decodes
 - R R E E D (RED, 0.7)
 - R R - E D (RED, 0.68)
 - R R E E E D (RED, 0.69)
 - T T E E E D (TED, 0.71)
 - TT-EED (TED, 0.3)
 - TT--ED (TED, 0.29)
- A greedy decode picks the most likely output: TED
- A decode that considers the sum of all alignments of the same final output will select RED
- Which is more reasonable?
- And yet, remarkably, greedy decoding can be surprisingly effective, when using decoding with blanks



Figure 1. Framewise and CTC networks classifying a speech signal. The shaded lines are the output activations, corresponding to the probabilities of observing phonemes at particular times. The CTC network predicts only the sequence of phonemes (typically as a series of spikes, separated by 'blanks', or null predictions), while the framewise network attempts to align them with the manual segmentation (vertical lines). The framewise network receives an error for misaligning the segment boundaries, even if it predicts the correct phoneme (e.g. 'dh'). When one phoneme always occurs beside another (e.g. the closure 'dcl' with the stop 'd'), CTC tends to predict them together in a double spike. The choice of labelling can be read directly from the CTC outputs (follow the spikes), whereas the predictions of the framewise network must be post-processed before use.

- Ref: Graves
- Symbol outputs peak at the ends of the sounds
 - Typical output: - R - E - -D
 - Model output naturally eliminates alignment ambiguities
- But this is still suboptimal..

Actual objective of decoding

- Want to find most likely order-aligned symbol sequence
 - RED
 - What greedy decode finds: most likely time synchronous symbol sequence
 - -/R//R/--/EH//EH//D/
 - Which must be compressed
- Find the order-aligned symbol sequence $S = S_0, ..., S_{K-1}$, given an input $X = X_0, ..., X_{T-1}$, that is most likely

$$= \operatorname*{argmax}_{S} P(S_0, \dots, S_{K-1} | \mathbf{X})$$

Recall: The forward probability $\alpha(t, r)$



$$\alpha_{S_0..S_{K-1}}(T-1,K-1) = P(S_0..S_{K-1}|\mathbf{X})$$

• The probability of the entire symbol sequence is the alpha at the bottom right node

Actual decoding objective

• Find the most likely (asynchronous) symbol sequence

$$\widehat{\mathbf{S}} = \operatorname*{argmax}_{\mathbf{S}} \alpha_{\mathbf{S}}(S_{K-1}, T-1)$$

- Unfortunately, explicit computation of this will require evaluate of an exponential number of symbol sequences
- Solution: Organize all possible symbol sequences as a (semi)tree

Hypothesis semi-tree



- The semi tree of hypotheses (assuming only 3 symbols in the vocabulary)
- Every symbol connects to every symbol other than itself
 - It also connects to a blank, which connects to every symbol including itself
- The simple structure repeats recursively
- Each node represents a unique (partial) symbol sequence!

The decoding graph for the tree



- The figure to the left is the tree, drawn in a vertical line
- The graph is just the tree unrolled over time
 - For a vocabulary of V symbols, every node connects out to V other nodes at the next time
- Every node in the graph represents a unique symbol sequence

The decoding graph for the tree



- The forward score $\alpha(r, T)$ at the final time represents the full forward score for a unique symbol sequence (including sequences terminating in blanks)
- Select the symbol sequence with the largest alpha
 - Sequences may have two alphas, one for the sequence itself, one for the sequence followed by a blank
 - Add the alphas before selecting the most likely sequence



- This is the "theoretically correct" CTC decoder
- In practice, the graph gets exponentially large very quickly
- To prevent this pruning strategies are employed to keep the graph (and computation) manageable
 - This may cause suboptimal decodes, however
 - The fact that CTC scores peak at symbol terminations minimizes the damage due to pruning

Beamsearch Pseudocode Notes

- Retaining separate lists of paths and pathscores for paths terminating in blanks, and those terminating in valid symbols
 - Since blanks are special
 - Do not explicitly represent blanks in the partial decode strings
- Pseudocode takes liberties (particularly w.r.t null strings)
 - I.e. you must be careful if you convert this to code
- Key
 - **PathScore** : array of scores for paths ending with symbols
 - **BlankPathScore** : array of scores for paths ending with blanks
 - **SymbolSet** : A list of symbols *not* including the blank

```
Global PathScore = [], BlankPathScore = []
# First time instant: Initialize paths with each of the symbols,
# including blank, using score at time t=1
NewPathsWithTerminalBlank, NewPathsWithTerminalSymbol, NewBlankPathScore, NewPathScore =
                           InitializePaths(SymbolSet, y[:,0])
# Subsequent time steps
for t = 1:T
    # Prune the collection down to the BeamWidth
    PathsWithTerminalBlank, PathsWithTerminalSymbol, BlankPathScore, PathScore =
                   Prune (NewPathsWithTerminalBlank, NewPathsWithTerminalSymbol,
                                             NewBlankPathScore, NewPathScore, BeamWidth)
    # First extend paths by a blank
    NewPathsWithTerminalBlank, NewBlankPathScore = ExtendWithBlank(PathsWithTerminalBlank,
                                                                  PathsWithTerminalSymbol, y[:,t])
    # Next extend paths by a symbol
    NewPathsWithTerminalSymbol, NewPathScore = ExtendWithSymbol(PathsWithTerminalBlank,
                                                         PathsWithTerminalSymbol, SymbolSet, y[:,t])
```

end

Merge identical paths differing only by the final blank

Pick best path
BestPath = argmax(FinalPathScore) # Find the path with the best score

Global PathScore = [], BlankPathScore = []

First time instant: Initialize paths with each of the symbols, # including blank, using score at time t=1 NewPathsWithTerminalBlank, NewPathsWithTerminalSymbol, NewBlankPathScore, NewPathScore = InitializePaths(SymbolSet, y[:,0])





)

```
Global PathScore = [], BlankPathScore = []
```

First time instant: Initialize paths with each of the symbols, # including blank, using score at time t=1 NewPathsWithTerminalBlank, NewPathsWithTerminalSymbol, NewBlankPathScore, NewPathScore = InitializePaths(SymbolSet, y[:,0])

Subsequent time steps

for t = 1:T

Prune the collection down to the BeamWidth

```
PathsWithTerminalBlank, PathsWithTerminalSymbol, BlankPathScore, PathScore =

Prune (NewPathsWithTerminalBlank, NewPathsWithTerminalSymbol,

NewBlankPathScore, NewPathScore, BeamWidth)
```





)

```
Global PathScore = [], BlankPathScore = []
```

First time instant: Initialize paths with each of the symbols, # including blank, using score at time t=1

NewPathsWithTerminalBlank, NewPathsWithTerminalSymbol, NewBlankPathScore, NewPathScore =
InitializePaths(SymbolSet, y[:,0])

Subsequent time steps

for t = 1:T

Prune the collection down to the BeamWidth

PathsWithTerminalBlank, PathsWithTerminalSymbol, BlankPathScore, PathScore = **Prune**(NewPathsWithTerminalBlank, NewPathsWithTerminalSymbol,

NewBlankPathScore, NewPathScore, BeamWidth)





)

```
Global PathScore = [], BlankPathScore = []
```

First time instant: Initialize paths with each of the symbols, # including blank, using score at time t=1

NewPathsWithTerminalBlank, NewPathsWithTerminalSymbol, NewBlankPathScore, NewPathScore =
InitializePaths(SymbolSet, y[:,0])

Subsequent time steps

for t = 1:T

Prune the collection down to the BeamWidth

NewBlankPathScore, NewPathScore, BeamWidth)

First extend paths by a blank NewPathsWithTerminalBlank, NewBlankPathScore = ExtendWithBlank(PathsWithTerminalBlank, PathsWithTerminalSymbol, y[:,t])



BEAM SEARCH InitializePaths: FIRST TIME INSTANT

```
function InitializePaths(SymbolSet, y)
```

```
InitialBlankPathScore = [], InitialPathScore = []
# First push the blank into a path-ending-with-blank stack. No symbol has been invoked yet
path = null
InitialBlankPathScore[path] = y[blank] # Score of blank at t=1
InitialPathsWithFinalBlank = {path}
```

Push rest of the symbols into a path-ending-with-symbol stack

```
InitialPathsWithFinalSymbol = {}
for c in SymbolSet # This is the entire symbol set, without the blank
    path = c
    InitialPathScore[path] = y[c] # Score of symbol c at t=1
    InitialPathsWithFinalSymbol += path # Set addition
```

end

```
return InitialPathsWithFinalBlank, InitialPathsWithFinalSymbol,
InitialBlankPathScore, InitialPathScore
```



BEAM SEARCH: Extending with blanks

Global PathScore, BlankPathScore

```
function ExtendWithBlank (PathsWithTerminalBlank, PathsWithTerminalSymbol, y)
   UpdatedPathsWithTerminalBlank = {}
   UpdatedBlankPathScore = []
    # First work on paths with terminal blanks
    #(This represents transitions along horizontal trellis edges for blanks)
    for path in PathsWithTerminalBlank:
        # Repeating a blank doesn't change the symbol sequence
       UpdatedPathsWithTerminalBlank += path # Set addition
       UpdatedBlankPathScore[path] = BlankPathScore[path]*y[blank]
    end
    # Then extend paths with terminal symbols by blanks
    for path in PathsWithTerminalSymbol:
        # If there is already an equivalent string in UpdatesPathsWithTerminalBlank
       # simply add the score. If not create a new entry
        if path in UpdatedPathsWithTerminalBlank
            UpdatedBlankPathScore[path] += Pathscore[path]* y[blank]
        else
            UpdatedPathsWithTerminalBlank += path # Set addition
            UpdatedBlankPathScore[path] = PathScore[path] * y[blank]
        end
    end
    return UpdatedPathsWithTerminalBlank,
          UpdatedBlankPathScore
```

103

BEAM SEARCH: Extending with symbols

```
Global PathScore, BlankPathScore
```

```
function ExtendWithSymbol (PathsWithTerminalBlank, PathsWithTerminalSymbol, SymbolSet, y)
   UpdatedPathsWithTerminalSymbol = {}
   UpdatedPathScore = []
   # First extend the paths terminating in blanks. This will always create a new sequence
   for path in PathsWithTerminalBlank:
       for c in SymbolSet: # SymbolSet does not include blanks
            newpath = path + c # Concatenation
            UpdatedPathsWithTerminalSymbol += newpath # Set addition
           UpdatedPathScore[newpath] = BlankPathScore[path] * y(c)
       end
   end
   # Next work on paths with terminal symbols
   for path in PathsWithTerminalSymbol:
        # Extend the path with every symbol other than blank
       for c in SymbolSet: # SymbolSet does not include blanks
            newpath = (c == path[end]) ? path : path + c # Horizontal transitions don't extend the sequence
           if newpath in UpdatedPathsWithTerminalSymbol: # Already in list, merge paths
               UpdatedPathScore[newpath] += PathScore[path] * y[c]
            else # Create new path
               UpdatedPathsWithTerminalSymbol += newpath # Set addition
               UpdatedPathScore[newpath] = PathScore[path] * y[c]
           end
       end
   end
```



BEAM SEARCH: Pruning low-scoring entries

Global PathScore, BlankPathScore

```
function Prune (PathsWithTerminalBlank, PathsWithTerminalSymbol, BlankPathScore, PathScore, BeamWidth)
   PrunedBlankPathScore = []
   PrunedPathScore = []
   # First gather all the relevant scores
   i = 1
   for p in PathsWithTerminalBlank
        scorelist[i] = BlankPathScore[p]
        i++
   end
   for p in PathsWithTerminalSymbol
        scorelist[i] = PathScore[p]
        i++
   end
   # Sort and find cutoff score that retains exactly BeamWidth paths
   sort(scorelist) # In decreasing order
   cutoff = BeamWidth < length(scorelist) ? scorelist[BeamWidth] : scorelist[end]</pre>
   PrunedPathsWithTerminalBlank = {}
   for p in PathsWithTerminalBlank
        if BlankPathScore[p] >= cutoff
            PrunedPathsWithTerminalBlank += p # Set addition
            PrunedBlankPathScore[p] = BlankPathScore[p]
        end
   end
   PrunedPathsWithTerminalSymbol = { }
   for p in PathsWithTerminalSymbol
        if PathScore[p] >= cutoff
            PrunedPathsWithTerminalSymbol += p # Set addition
            PrunedPathScore[p] = PathScore[p]
        end
```

end

return PrunedPathsWithTerminalBlank, PrunedPathsWithTerminalSymbol, PrunedBlankPathScore, PrunedPathScore

BEAM SEARCH: Merging final paths

Note : not using global variable here

```
# All paths with terminal symbols will remain
MergedPaths = PathsWithTerminalSymbol
FinalPathScore = PathScore
```

```
# Paths with terminal blanks will contribute scores to existing identical paths from
# PathsWithTerminalSymbol if present, or be included in the final set, otherwise
```

```
for p in PathsWithTerminalBlank
```

```
if p in MergedPaths
    FinalPathScore[p] += BlankPathScore[p]
```

else

```
MergedPaths += p # Set addition
FinalPathScore[p] = BlankPathScore[p]
```

end

end

```
return MergedPaths, FinalPathScore
```

Story so far: CTC models

- Sequence-to-sequence networks which irregularly produce output symbols can be trained by
 - Iteratively aligning the target output to the input and time-synchronous training
 - Optimizing the expected error over *all* possible alignments: CTC training
- Distinct repetition of symbols can be disambiguated from repetitions representing the extended output of a single symbol by the introduction of blanks
- Decoding the models can be performed by
 - Best-path decoding, i.e. Viterbi decoding
 - Optimal CTC decoding based on the application of the forward algorithm to a tree-structured representation of all possible output strings

Most common CTC applications

- Speech recognition
 - Speech in, phoneme sequence out
 - Speech in, character sequence (spelling out)
- Handwriting recognition


- Recurrent neural networks (with LSTMs) can be used to perform speech recognition
 - Input: Sequences of audio feature vectors
 - Output: Phonetic label of each vector

Speech recognition using Recurrent Nets



• Alternative: Directly output phoneme, character or word sequence

Next up: Attention models